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Markov Chain Monte Carlo and Its application for Travelling Salesman Problem

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**1. Theory**

**1.1 Markov Chains:**

Markov chains are random sequence of finite states and the transition probability for switching states is totally depends on the current states only.

Let us suppose we have a set of states, S = {s1,s2,. . . . . sr}. The process starts in one of these states and moves successively from one state to another. Each move is called a step. If the chain is currently in state si, then it moves to state sj at the next step with a probability denoted by pij , and this probability does not depend upon which states the chain was in before the current state.

The Probabilities are called Transition Probabilities and the matrix containing all such probabilities is called Transition Matrix. The process can remain in the state it is in, and this occurs with probability pii. An initial probability distribution, defined on S, specifies the starting state. Usually this is done by specifying a particular state as the starting state.

A sequence {Xn: n>=0} of random variables taking values in a countable state space S is called a Marcov Chain if,

P {Xn =αn|Xn-1=αn-1, Xn-2=αn-2……X0=α0} = P {Xn=αn | Xn-1=αn-1}

I.e. the probability that the chain will be in a certain state αn at time n given all its past history depends only on its previous state at time n-1.

Usually we also impose the condition of homogeneity

P {Xn+1=j | Xn=i} = P {X1=j | X0=i}

**1.2 Markov Chain Monte Carlo:**

Markov chain Monte Carlo (MCMC) methods are a class of algorithms for sampling from probability distributions based on constructing a Markov chain that has the desired distribution as its equilibrium distribution. It is generally used when the random vector to be simulated is composed of a large number of independent variables to generate Markov chain whose equilibrium distribution is equal to required distribution.

**1.3 Simulated Annealing:**

Simulated Annealing is a generic probabilistic meta-algorithm used to find an approximate solution to global optimization problems. It is inspired by annealing in metallurgy which is a technique of controlled cooling of material to reduce defects. The simulated annealing algorithm starts with a random solution. Each iteration forms a random nearby solution. If this solution is a better solution, it will replace the current solution. If it is a worse solution, it may be chosen to replace the current solution with a probability that depends on the temperature parameter. As the algorithm progresses, the temperature parameter decreases, giving worse solutions a lesser chance of replacing the current solution. Allowing worse solutions at the beginning helps to avoid converging to a local minimum rather than the global minimum.

**1.4 Travelling Salesman Problem:**

The travelling salesman problem (TSP) asks the following question: Given a list of cities and the distances between each pair of cities, what is the shortest possible route that visits each city exactly once and returns to the origin city? It is an NP-hard problem and is used as a benchmark for many optimization methods.

The TSP has several applications even in its purest formulation, such as planning, logistics, and the manufacture of microchips. Slightly modified, it appears as a sub-problem in many areas, such as DNA sequencing. In these applications, the concept city represents, for example, customers, soldering points, or DNA fragments, and the concept distance represents travelling times or cost, or a similarity measure between DNA fragments.

**2. Project Description:**

**2.1 Simulated Annealing (SA) for Travelling Salesman problem**

We took 20 cities for the simulation purposes so our state space is set of unique the permutation of those 20 cities and these permutation is called tour. The size of our state space is 20! (Finite).

The algorithm stats with a tour which is just a random sequence of 20 cities and will improve gradually along the time. Our algorithm’s next state which will be compared and evaluated with current state is obtained by randomly switching two cities in the current tour.

We took total distance that needs to be travelled for certain tour to be energy of the state and our goal is to minimize the energy of the state.

Temperature is the key to SA's ability to find global optimums, and what differentiates it from simple hill-climbing algorithms, is that at each iteration it has a (decreasing) probability of moving to more energy having states. It always moves to more less energy having states but, when next state is more energy having state, it will move to next state when is grater then random (0,1) variable, where T is temperature. Introduction of random variable randomizes the next state which is the key differentiator of simulated annealing from simple annealing.

Theoretical analysis of SA indicates a cooling rate will determine the probability of escaping from local optimums and eventually finding the global optimum because as time passes based on cooling rate with the time temperature decreases reducing the probability of moving to the more energy having states. In our simulation the temperature is **104** and cooling rate is **0.01%.**

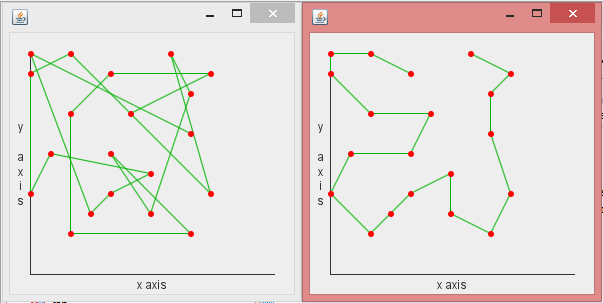
We have divided our project following major sections.

* Generating and managing cities
* Tour management
* Simulated Annealing
* Plotting the results

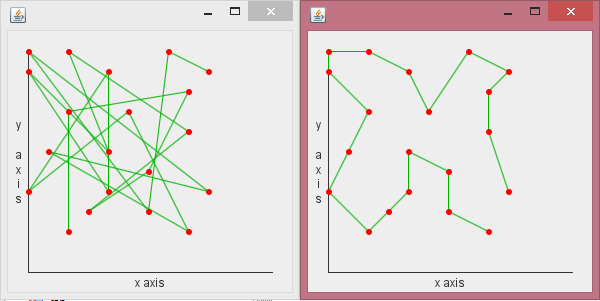
**3. Results:**

Here, we have displayed results of two different simulations. The best possible solution is a tour having a distance equal to 863 which we got by running the simulation for more than 100 times.

Figure has initial tour and the best possible tour that we got after first simulation.

 Initial tour: 2452 Final Tour: 909

This tour shows same result for another simulation.



Initial: 2371 Final: 863

Here is the data of sample 10 simulation that we ran.

|  |  |  |
| --- | --- | --- |
| **Simulation No** | **Initial Distance of the Tour** | **Best tour distance** |
| 1 | 2616 | 863 |
| 2 | 2467 | 863 |
| 3 | 1840 | 926 |
| 4 | 2011 | 881 |
| 5 | 2346 | 916 |
| 6 | 2064 | 863 |
| 7 | 2579 | 903 |
| 8 | 2583 | 947 |
| 9 | 2371 | 987 |
| 10 | 2452 | 909 |

**4. Work Distribution:**

1. Nikunj Amipara (201001199) – Research on Simulated Annealing, Coding the Simulated Annealing part of code.
2. Parth Shah (201001200) – Research on Markov Chains, Coding City and Tour part of code
3. Om Thakkar (201001203) – Research on Monte Carlo and applications of MCMC, Coding graphical part of code

**5. References:**

* Markov Chain Monte Carlo Simulation Made Simple by Alastair Smith
* Introduction to Markov Chain Monte Carlo by Charles J. Geyer
* Markov Chain Monte Carlo Method and Its Application by Stephen P. Brooks